

SibylFS and ImpFS: formal specification and implementation for filesystems

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REMS workshop, Cambridge, 2016-05-27

What have we done? SibylFS

- ▶ SibylFS, a filesystem specification and test oracle (POSIX, Mac, Linux and BSD filesystems)
- ▶ Spec formal, mechanized, written in Lem (translated to HOL, Isabelle, OCaml)
- ▶ Test oracle can analyse traces of real-world filesystems
- ▶ Tested lots of combinations of libc/OS/filesystem

Amusing application

- ▶ We presented at SOSP'15
- ▶ Also at that venue was a paper on FSCQ, a filesystem verified in Coq

dsheets commented on 3 Oct 2015

On an empty file system built from [27c4078](#), running

```
#include <unistd.h>
#include <fcntl.h>

int main() {
    close(creat("spaaaaaaace"));
    truncate("spaaaaaaace", 10);

    return 0;
}
```

results in a file, `spaaaaaaace`, which contains the text `spaaaaaaace` instead of 10 null bytes. For applications which expect POSIX, OS X, Linux, or BSD behavior when using `truncate` to extend files, this unexpected data can result in incorrectly formatted files.



zeldovich commented on 3 Oct 2015

Owner

Thanks; I will take a look. Sounds like we might have gotten the spec wrong.

Structure of the specification

- This is relevant to what we are trying to do now...

Call to libc

|
V

SibylFS spec
(3 or 4 layers)

|
V

Dir_heap // this is essentially two maps

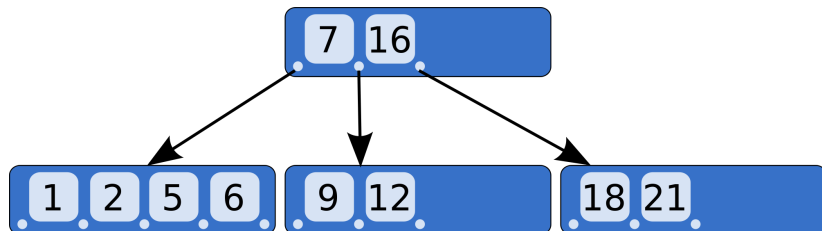
- Dir_heap is 2 maps, one from file ids to files, one from directory ids to directories

What are we doing now? ImpFS

- ▶ ImpFS, a verified filesystem implementation (from the start, this was what I really wanted to build)
- ▶ Based on B-trees

What are B-trees?

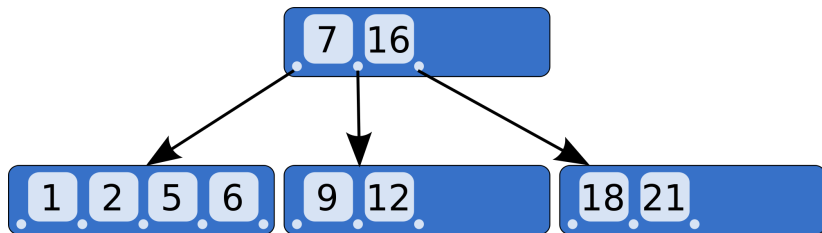
- ▶ A balanced tree-like datastructure, implementing a map:



- ▶ This shows the keys; in the leaves, we also have values associated with each key; the keys in the nodes allow you to navigate to a desired key in a leaf

Why are B-trees important in filesystems (and databases etc)?

- ▶ Nodes are mapped onto blocks
- ▶ Block reads are slow compared to everything else
- ▶ Tree structure minimizes the number of blocks that are traversed (read from disk) to locate a value for a given key
- ▶ N.B. two datastructures - the tree on disk, and the partial tree in memory (e.g. the stack of blocks leading to a desired leaf)
- ▶ Moreover, if we update a map (add/remove a key) we can share most of the blocks with the previous version of the map



How to get a verified filesystem?

- ▶ Take the SibylFS specification
- ▶ Remove non-determinism, so it becomes an executable in-memory filesystem
- ▶ Implement the two maps at the bottom of the spec layers, and interface to a block device, to get a “real” filesystem using block storage; B-trees are used to implement maps on top of a block device

Structure of ImpFS

Call to libc

|

V

SibylFS spec (determinised)

|

V

Dir_heap // two maps, implemented using B-trees

|

V

Block device

Of course, it is never that easy (challenges)

- ▶ Block layer (asynchronous; reorders writes; syncable; behaviour on crash; on-disk layout of data; pointers everywhere etc)
- ▶ Concurrency (we want an arbitrary number of B-trees concurrently using the block layer, ensuring all invariants are respected)
- ▶ **B-tree implementation and proof of correctness** taking account of block layer and concurrency requirements (I am most proud of this bit)
- ▶ Caching (want to avoid writes to block layer where possible)
- ▶ GC for unreachable blocks
- ▶ We want bounded memory usage (so have fixed size block caches etc)
- ▶ Performance
- ▶ “Semantic” optimizations
- ▶ Extra features (snapshotting etc)

ImpFS status

- ▶ Working on finishing the B-tree implementation and proof (should be an OCaml library and Isabelle theories in next few months)
- ▶ We hope for working filesystem prototype before October

Reuse

- ▶ B-tree proofs are very nice, and hopefully will form the basis for modern descriptions of this datastructure
- ▶ B-tree implementation (OCaml, Scala etc)
- ▶ ImpFS (in OCaml) could be used as a filesystem for Mirage, or for per-application private filesystems, or for filesystems on non-traditional block devices (network connections etc)
- ▶ B-tree can be used in many other applications (databases; per-application persistent kv maps etc)

End